

CS-523 Advanced topics on Privacy Enhancing Technologies

Censorship resistance

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Introduction

Censorship resistance

Course aim: learn **toolbox for privacy engineering**



toolbox
to enable free use of
digital communications



mechanisms
to evade censorship



attacks
that detect censorship
evasion

Application Layer

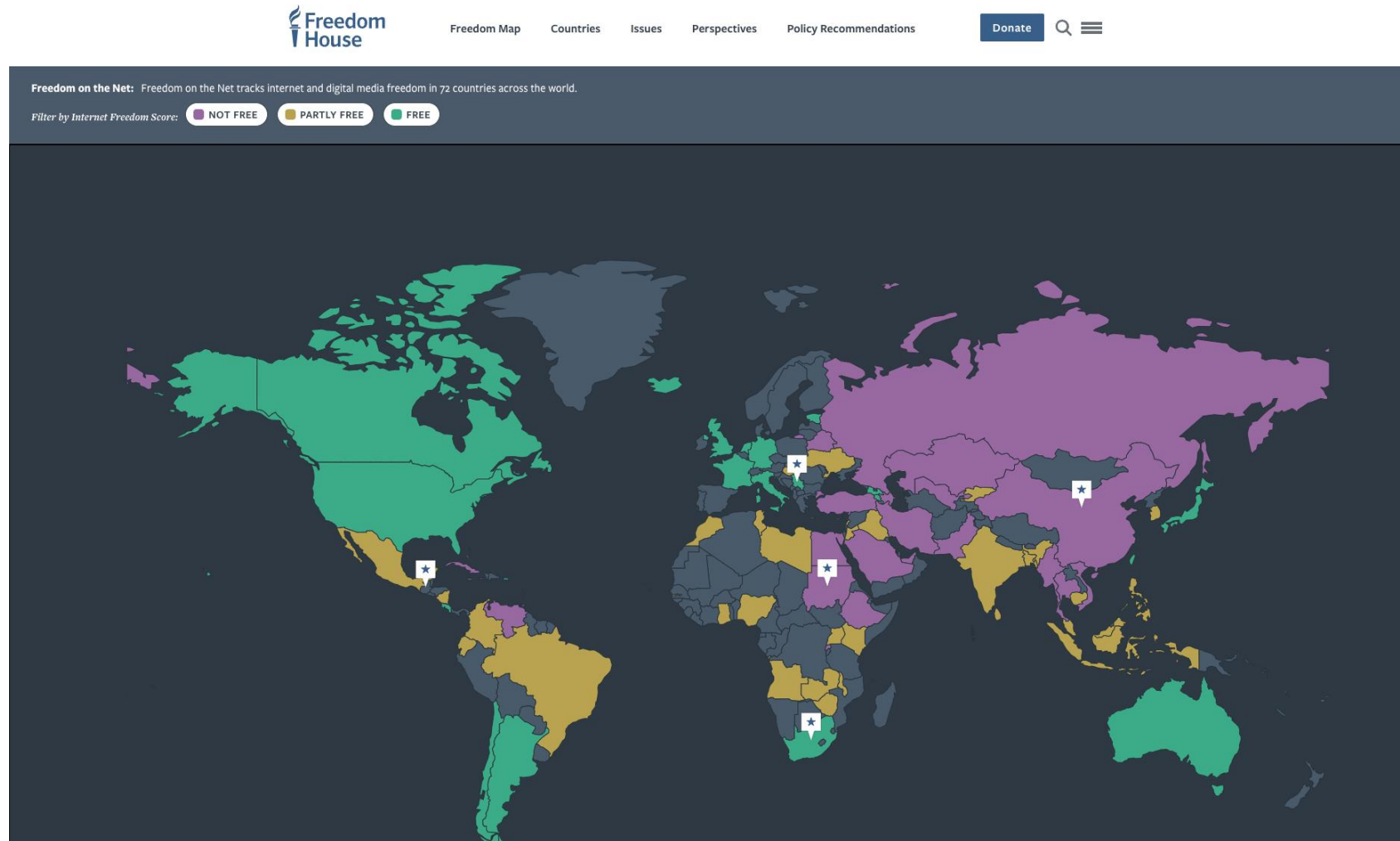
Network Layer

Goals

What should you learn today?

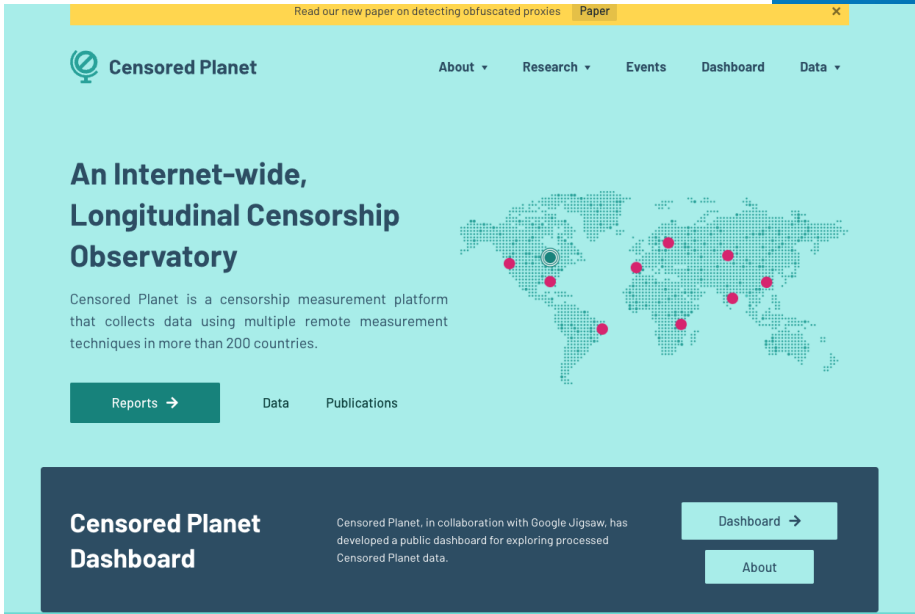
- Understand why **censorship resistance** is a privacy problem
- Understanding the **key elements** of censorship resistance
 - **Hide existence of communication**
 - **Enable communication**
- Which are the **most common pitfalls** of censorship resistance systems

Internet censorship is a global, evolving issue...



Internet censorship is a global, evolving issue...

How do we know?



<https://censoredplanet.org/censoredplanet>

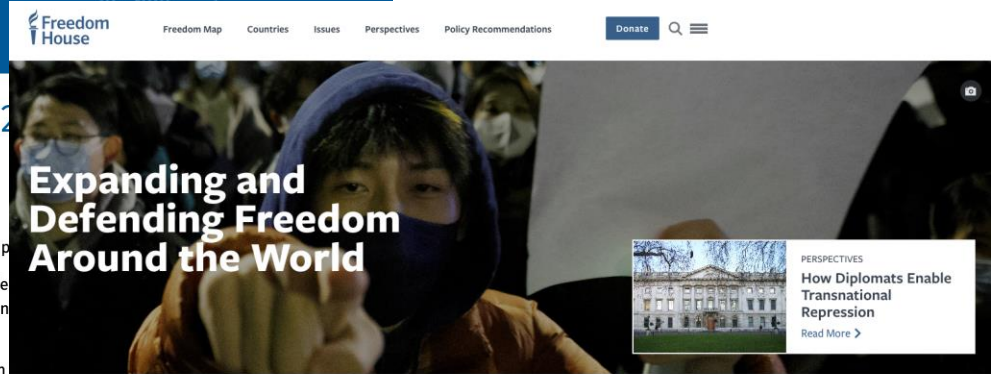


OONI Explorer is an open data resource on internet censorship. Since 2012, millions of network measurements have been collected from countries. OONI Explorer sheds light on internet censorship and interference worldwide.

To contribute to this open dataset, [install OONI Probe](#) and run

<https://explorer.ooni.org/>

<https://freedomhouse.org/>



Freedom House is founded on the core conviction that freedom flourishes in democratic nations where governments are accountable to their people.

Why censorship resistance?

- One of the goals of privacy technologies:
Self-determination

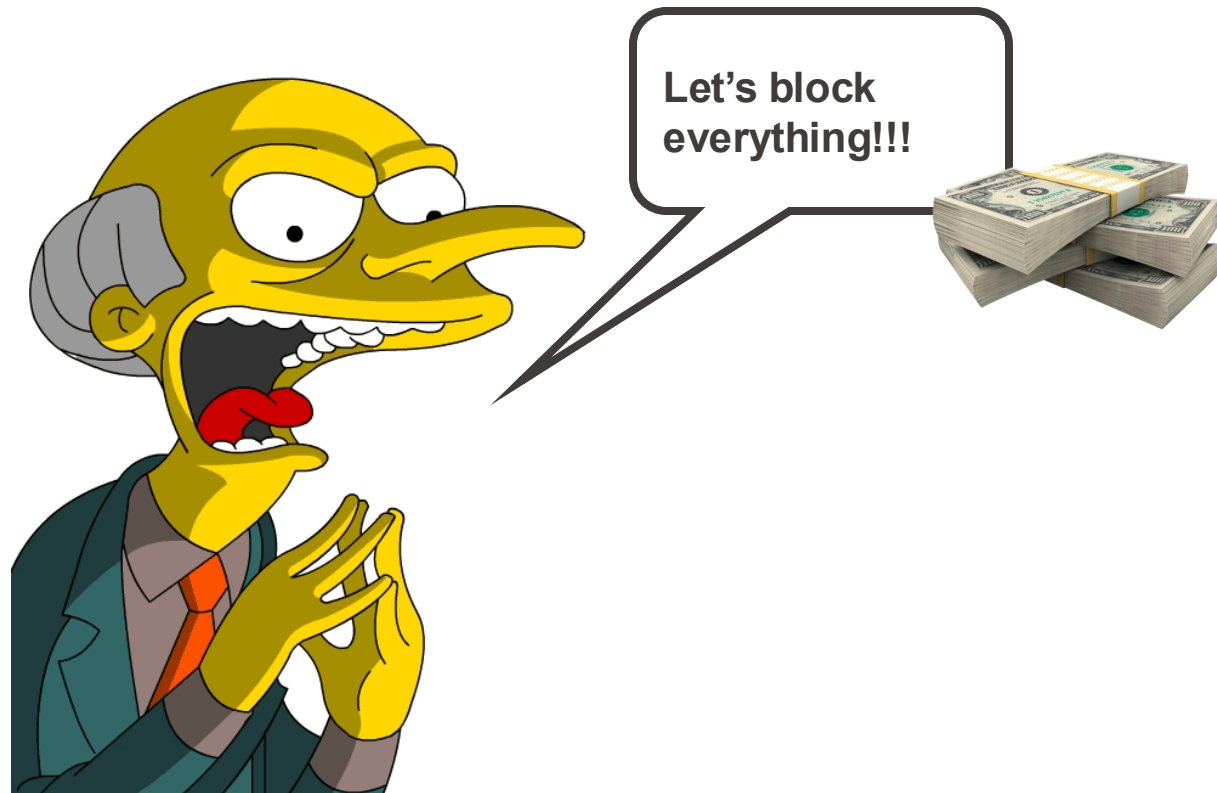
- Freedom of speech & freedom of information

Oh Jeez, where have I
seen some systems
that do this?

- Resisting Internet Censorship requires
 - re-routing : to avoid direct IP censorship
 - encryption : to avoid content-based censorship



Adversary's goal: prevent communication between two parties



Adversary's goal: prevent communication between two parties

An abstract model of censorship:



Step 1: Find the flow

Fingerprinting



Step 2: Prevent communication

Direct censor

Censorship

Step 1: Fingerprinting

Destination:

IP addresses, hosts, ports,...

Content:

protocol-strings, keywords, domains, http hosts, encrypted flows...

Flow properties:

length, inter-arrival times, bursts, ...

Protocol semantics:

protocol behavior (mostly active attacks)

Censorship



Step 2: Direct censor

Block destination:

Great Firewall of China



Degrade performance:

disrupt traffic, complicate access (soft form of censorship)

Corrupt routing:

BGP hijacking (disconnect part of the network)

DNS manipulation (redirect to censor or blackhole)



Corrupt flow content or semantics:

HTTP 404 not found

Forged RST packets

User-side/Publisher-side censorship:

local software/manual deletion

Goal of a censorship resistance system (CRS): unblockable communication between user and publisher*

Key components of CRS functionality:



Phase 1: Communication establishment
Get credentials



Phase 2: Conversation
Exchange information

Censorship resistance



Phase 1: Communication establishment

What: Obtain credentials or censorship resistance server addresses

Goal: **Easy** for users but **difficult** to censor

How: Hard to obtain/enumerate

- *High churn*: credentials/servers change continuously
- *Rate limit*: based on time, based on “space”, proof-of-work
- *Trust-based*: social graph, previous behavior, token,...

Active probing resistance:

- *Obfuscate aliveness*: only respond if correct sequence
- *Obfuscate service*: only respond with hidden service if correct sequence

Phase 2: Conversation

What: Actual communication

Goal: Avoid detection and blocking or modification of the conversation

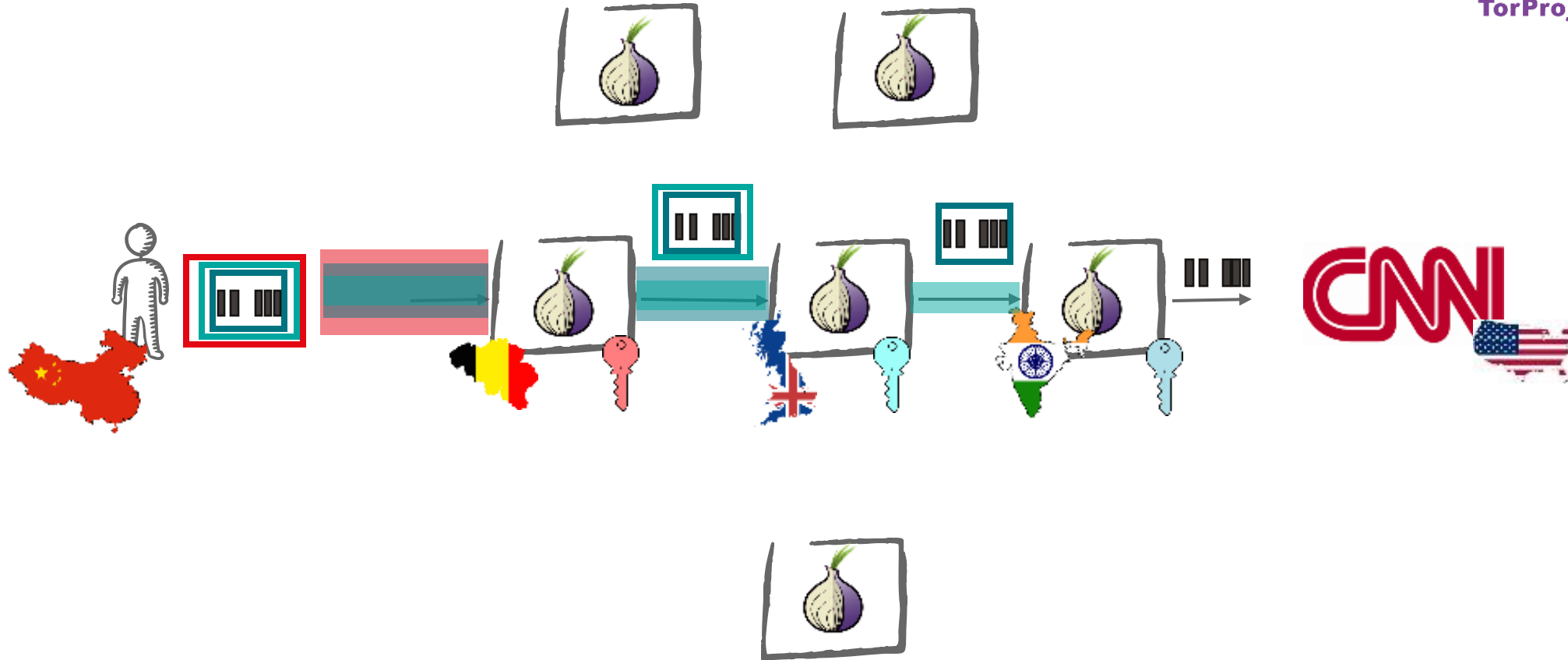
How: Destination obfuscation

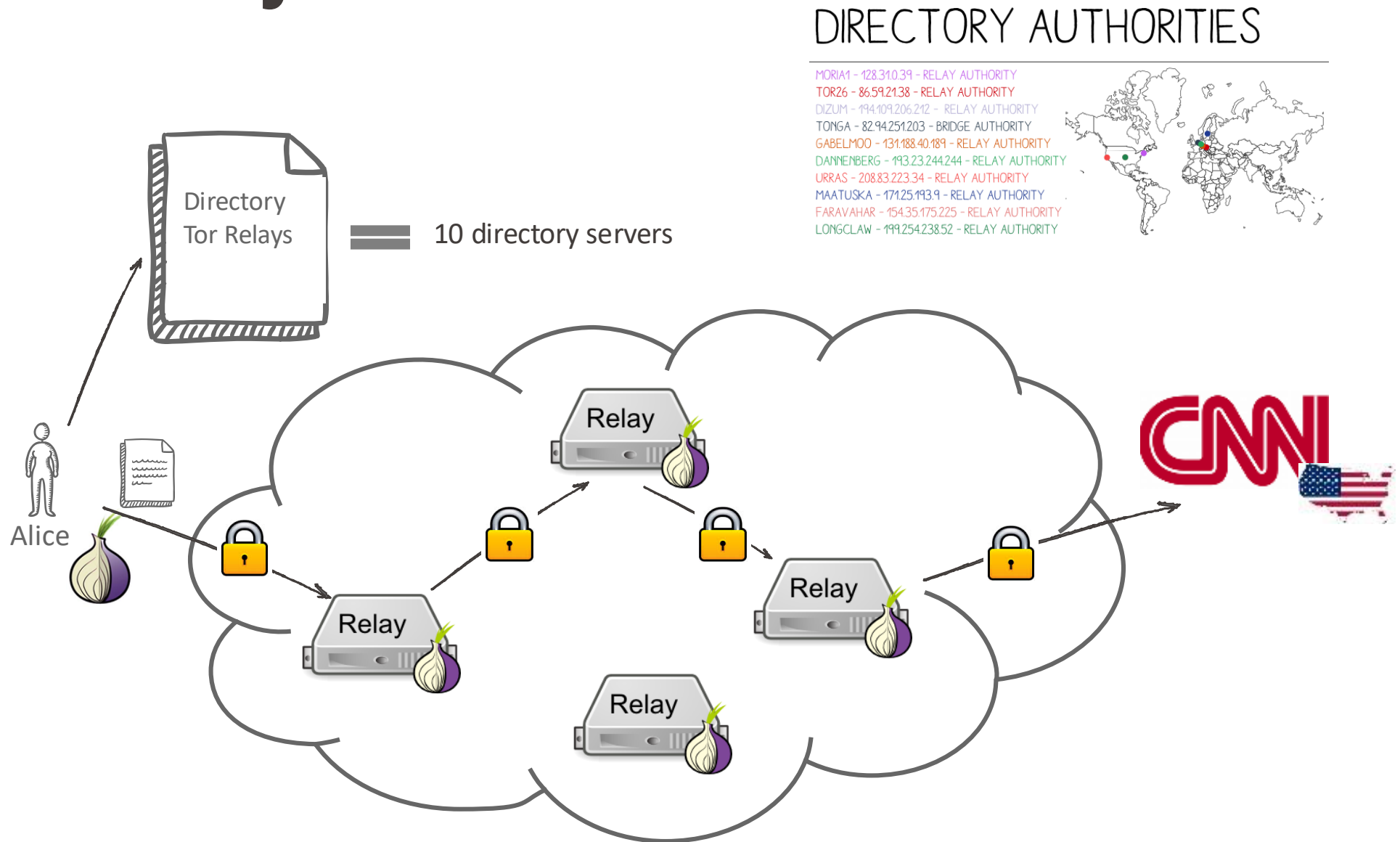
- **Proxy-based:** Tor
- **Decoy routing:** Telex, Cirripede,...

Content/flow obfuscation

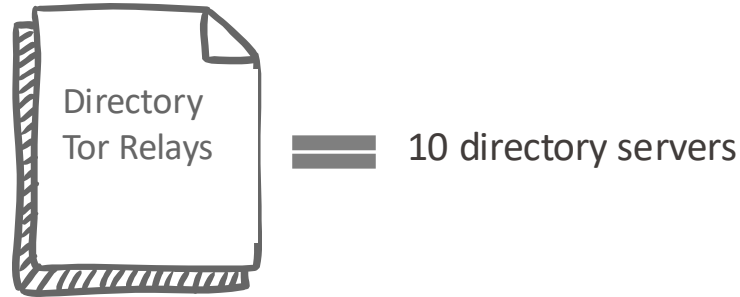
- **Mimicry:** look like whitelisted (or not blacklisted) ← increase cost of blocking
- **Tunneling:** tunnel traffic through unblocked application
- **Covert channel:** hide censored traffic on images, voice, emails,...

Tor as a CRS



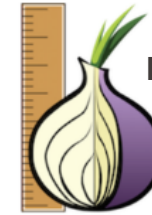


Tor directory authorities



DIRECTORY AUTHORITIES

MORIA1 - 128.310.39 - RELAY AUTHORITY
TOR26 - 86.59.21.38 - RELAY AUTHORITY
DIZUM - 194.109.206.212 - RELAY AUTHORITY
TONGA - 82.94.251.203 - BRIDGE AUTHORITY
GABELMOO - 131.188.40.189 - RELAY AUTHORITY
DANNENBERG - 193.23.244.244 - RELAY AUTHORITY
URRAS - 208.83.223.34 - RELAY AUTHORITY
MAATUSKA - 171.25.193.9 - RELAY AUTHORITY
FARAVAHAR - 154.35.175.225 - RELAY AUTHORITY
LONGCLAW - 199.254.238.52 - RELAY AUTHORITY



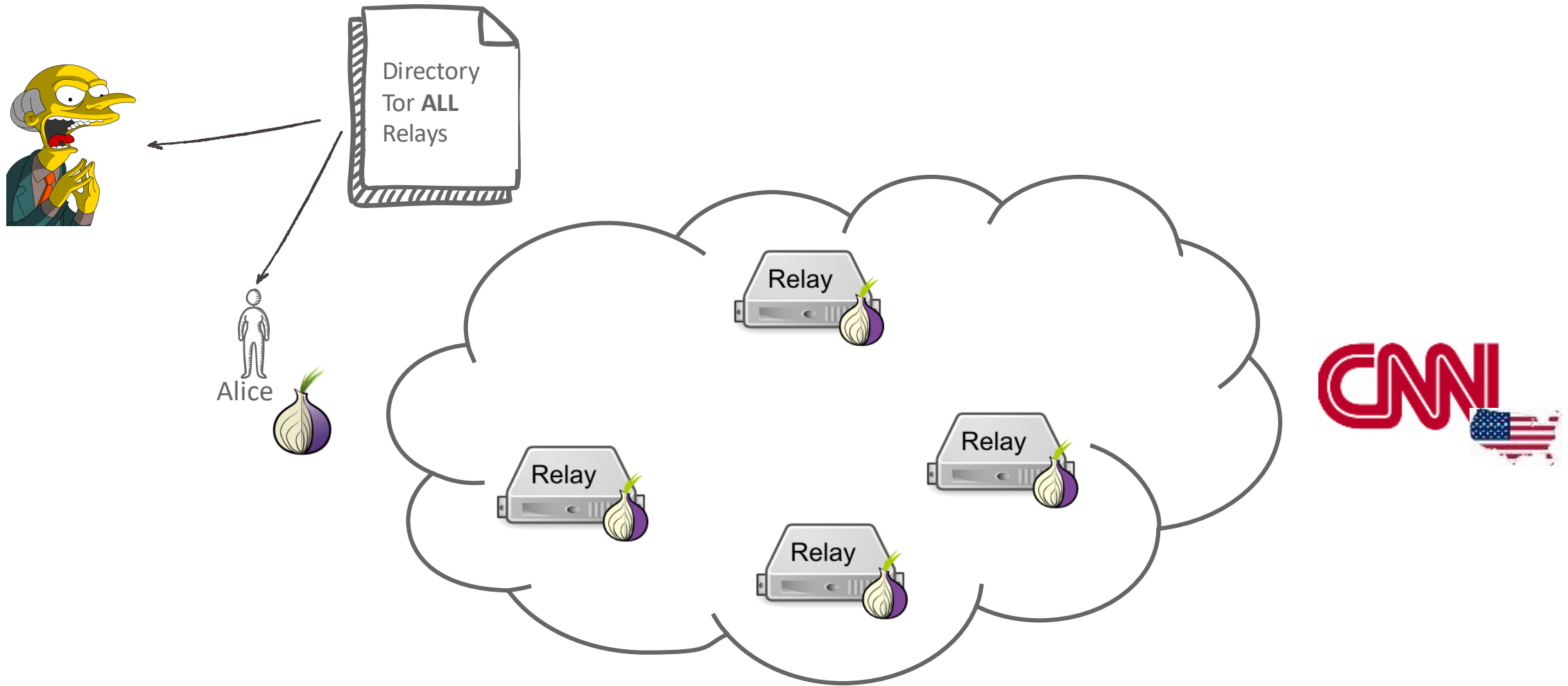
<https://metrics.torproject.org/collector.html>

TorMETRICS

Every hour:

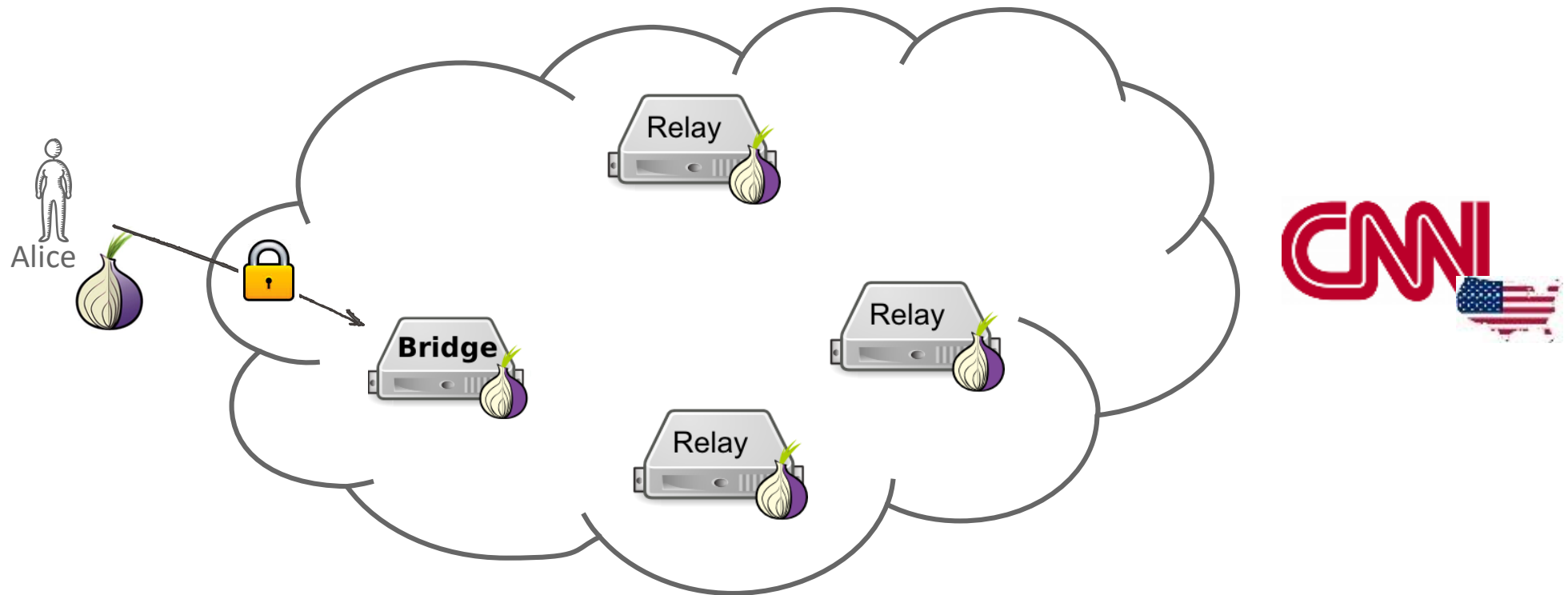
- Directory Authorities (DAs) compile a list of all known relays & flags & stuff
- DAs submit this “status-vote” to all the other authorities
- DAs combine parameters, sign and send to the other DA's
- There **should** be a majority agreeing on the data -> **consensus**
- **Consensus** published by each DA

Can an adversary block Tor?



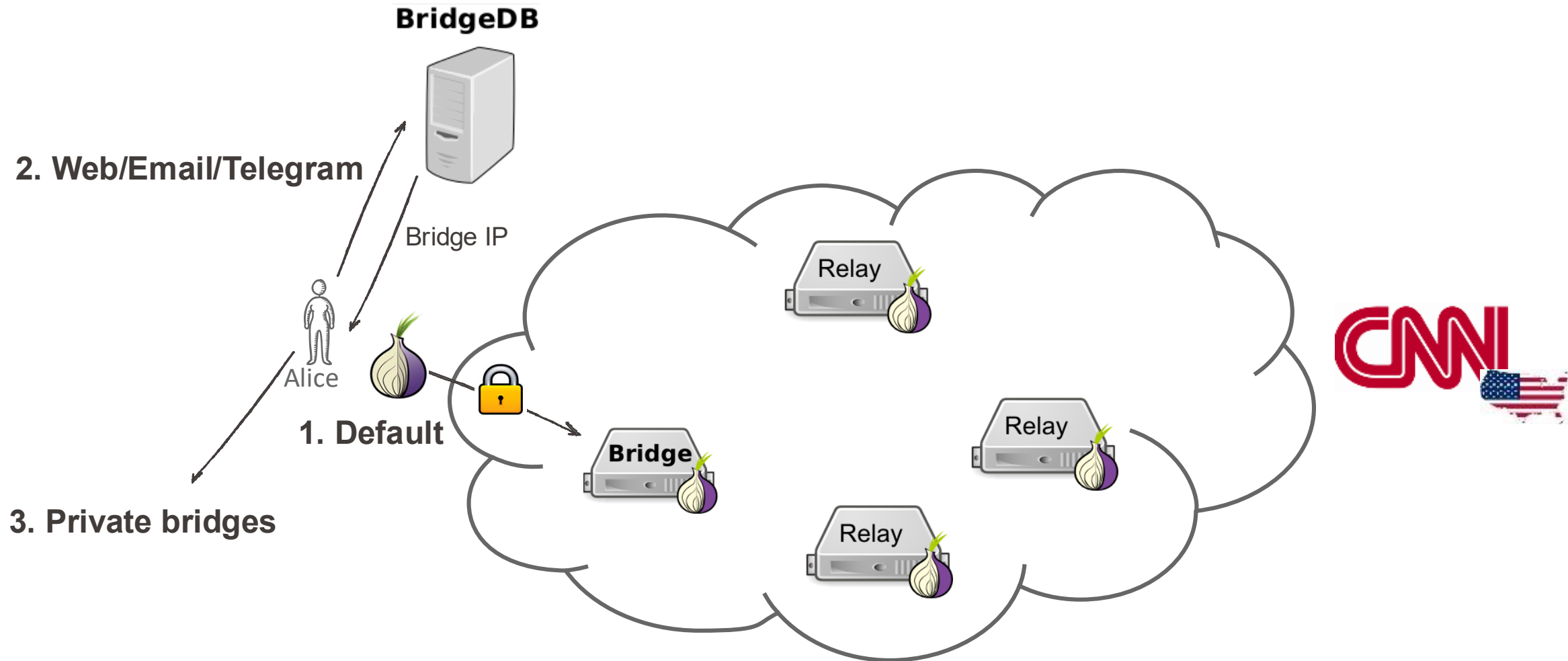
Tor bridges

Tor bridges: Onion routers whose IP is not publicly listed



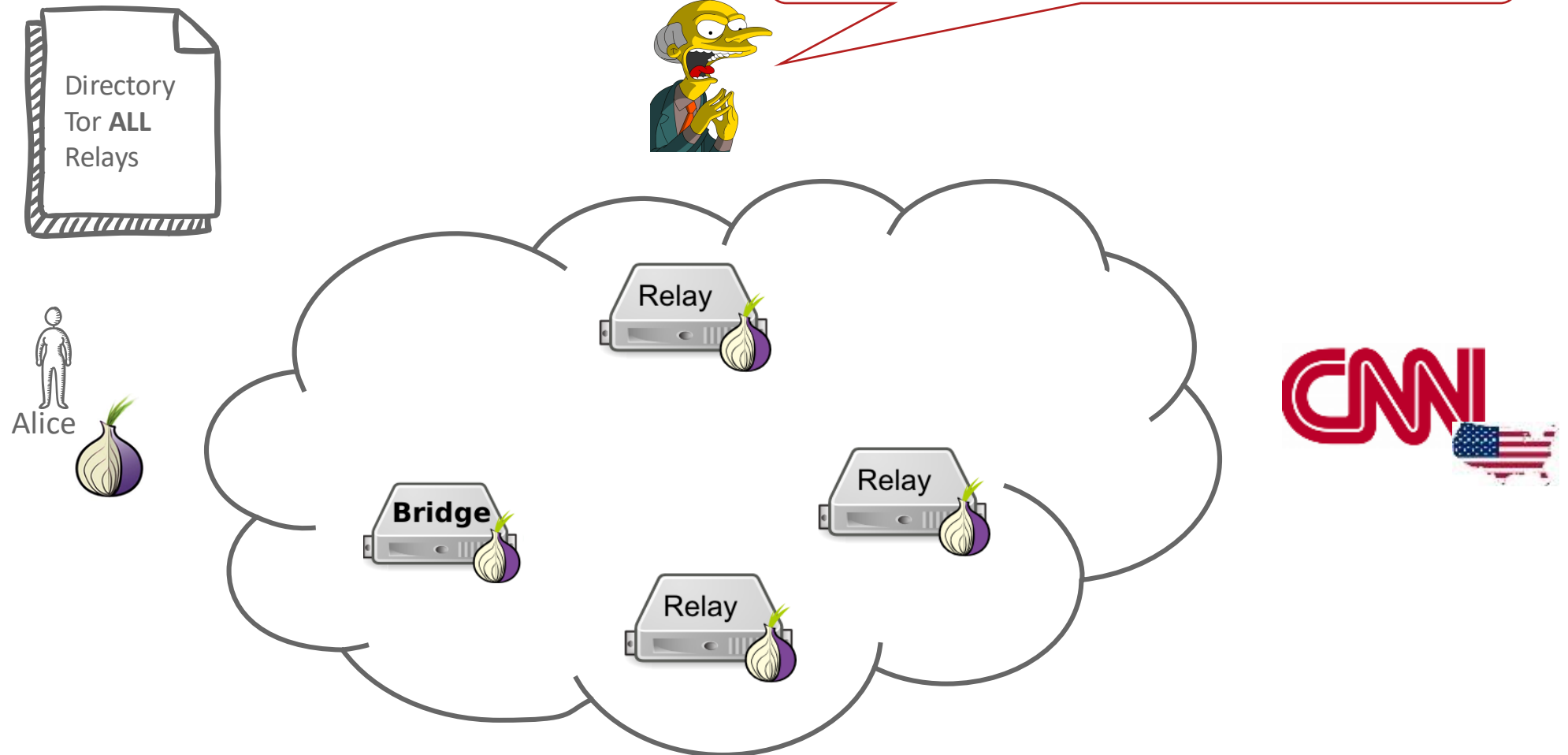
Tor bridges

How to find them



Tor bridges

Enumeration by censors



Tor bridges

Enumeration by censors

Two open Tor issues that censors can leverage to discover bridges

Issue 1: Vanilla Tor Certificates



- Vanilla Tor uses TLS handshake
- Easy to spot certificates
- It won't be fixed

Issue 2: Open Onion Routing Port



- Bridges have open OR Port with Vanilla Tor
- Even if they do not offer Vanilla Tor
- Difficult to fix

Pluggable transports



Currently there are four pluggable transports available, but more are being developed.

obfs4	obfs4 makes Tor traffic look random, and also prevents censors from finding bridges by Internet scanning. obfs4 bridges are less likely to be blocked than its predecessors, obfs3 bridges.
meek	meek transports make it look like you are browsing a major web site instead of using Tor. meek-azure makes it look like you are using a Microsoft web site.
Snowflake	Snowflake routes your connection through volunteer-operated proxies to make it look like you're placing a video call instead of using Tor.
WebTunnel	WebTunnel masks your Tor connection, making it appear as if you're accessing a website via HTTPS.

2025

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Content/flow obfuscation

- **Mimicry**: look like whitelisted (or not blacklisted)
- **Tunneling**: tunnel traffic through unblocked application
- **Covert channel**: hide censored traffic on images, voice, emails,...

Mimicry: Look like not blacklisted

ScrambleSuit

Key features

- Defense against active probing
- Pseudo-random payload
- Polymorphic

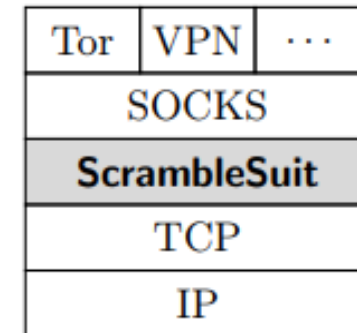


Figure 1:
ScrambleSuit's
protocol stack.

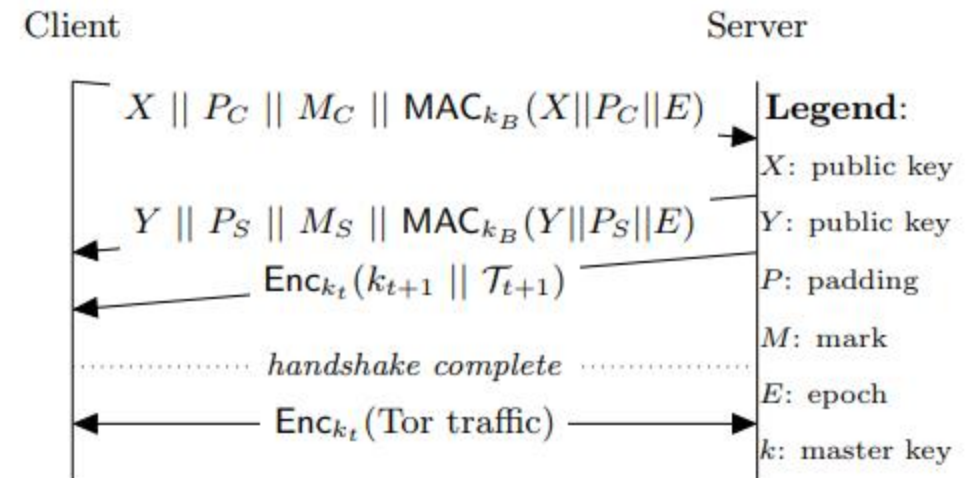
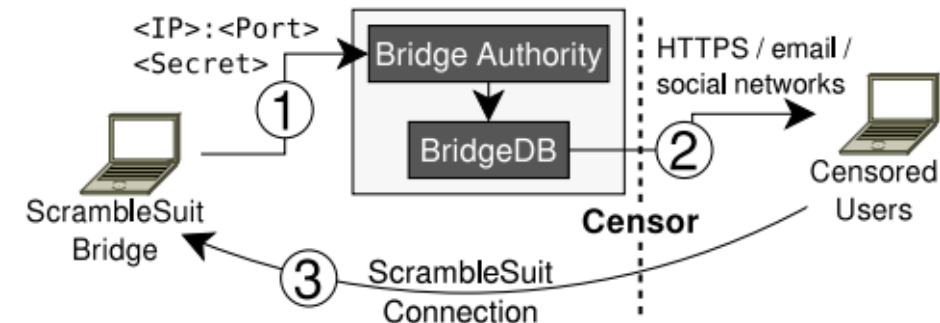
ScrambleSuit

Active attacks defense

Defense against active probing:

Protection against active probing attacks by requiring a *shared secret* between the client and the server.

This secret is communicated *out-of-band* via Tor's BridgeDB.



ScrambleSuit

Destroy patterns

(Lightweight*) Traffic analysis resistance through protocol polymorphism:

Every ScrambleSuit server generates its own and unique “protocol shape” by modifying:

- **packet lengths**
- **inter-arrival times**

How:

1. Generate a random seed shared between ScrambleSuit server and client.
2. Both sides use seed to generate two discrete probability distributions
3. Use distributions to shape traffic

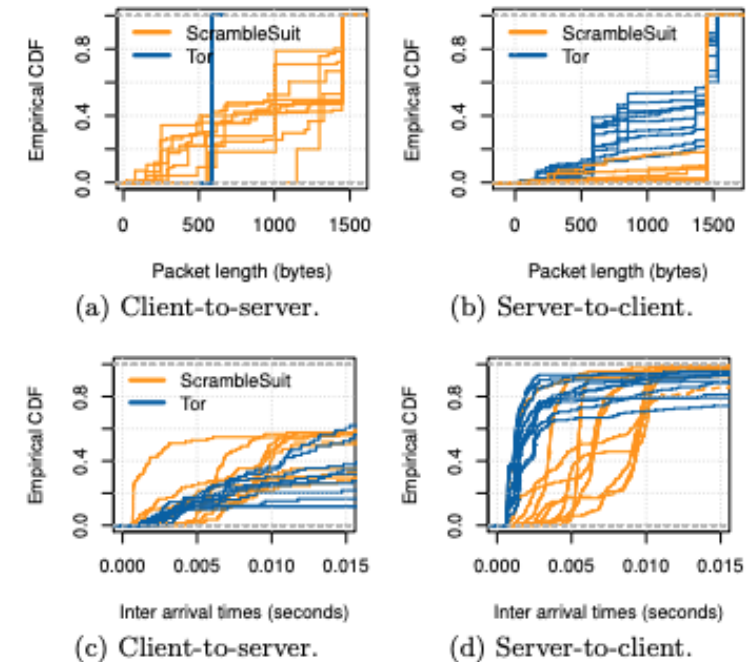


Figure 10: Tor’s and ScrambleSuit’s packet length distribution and inter-arrival times for both, client-to-server and server-to-client traffic.

ScrambleSuit

Destroy patterns

(Lightweight*) Traffic analysis resistance through protocol polymorphism:

Evaluation: How well does this mitigate?

Winter et al., 2013: “It is difficult to evaluate the effectiveness of our obfuscation techniques since ScrambleSuit **does not have a cover protocol to mimic**. Otherwise, our evaluation would simply investigate the similarity between our protocol and its cover protocol. Instead of measuring ScrambleSuit’s closeness to a mimicked protocol, we measure the deviation from its transported application, i.e., Tor. **Intuitively, higher deviation would imply better obfuscation.**”

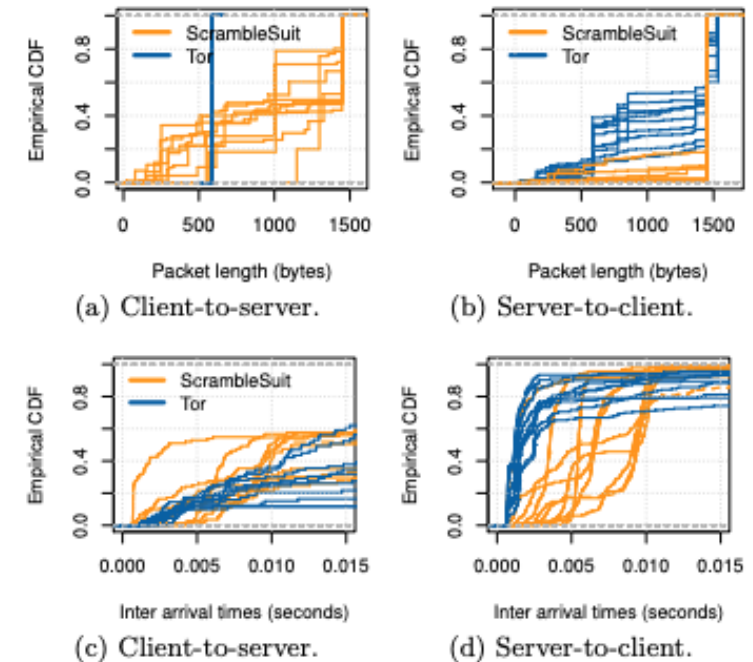


Figure 10: Tor’s and ScrambleSuit’s packet length distribution and inter-arrival times for both, client-to-server and server-to-client traffic.

ScrambleSuit

Destroy patterns

It looks like nothing...
... but nothing looks like it!

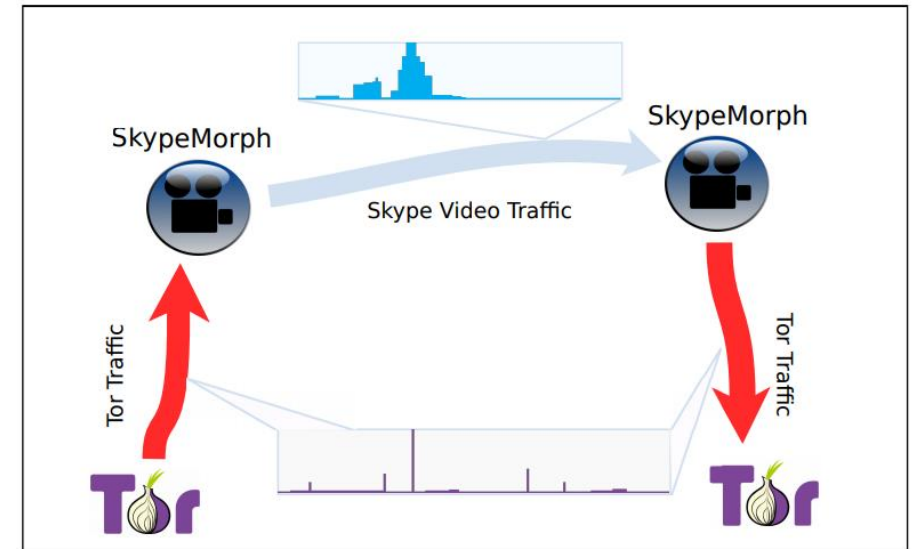


Mimicry: Look like whitelisted SkypeMorph

Goal: make it difficult for the censor to distinguish between the obfuscated bridge connections and whitelisted traffic using statistical comparisons of flow features

How: Tor clients obfuscate their messages to Tor bridge server in a widely used protocol over the Internet.

Skype video calls as target protocol



SkypeMorph Setup

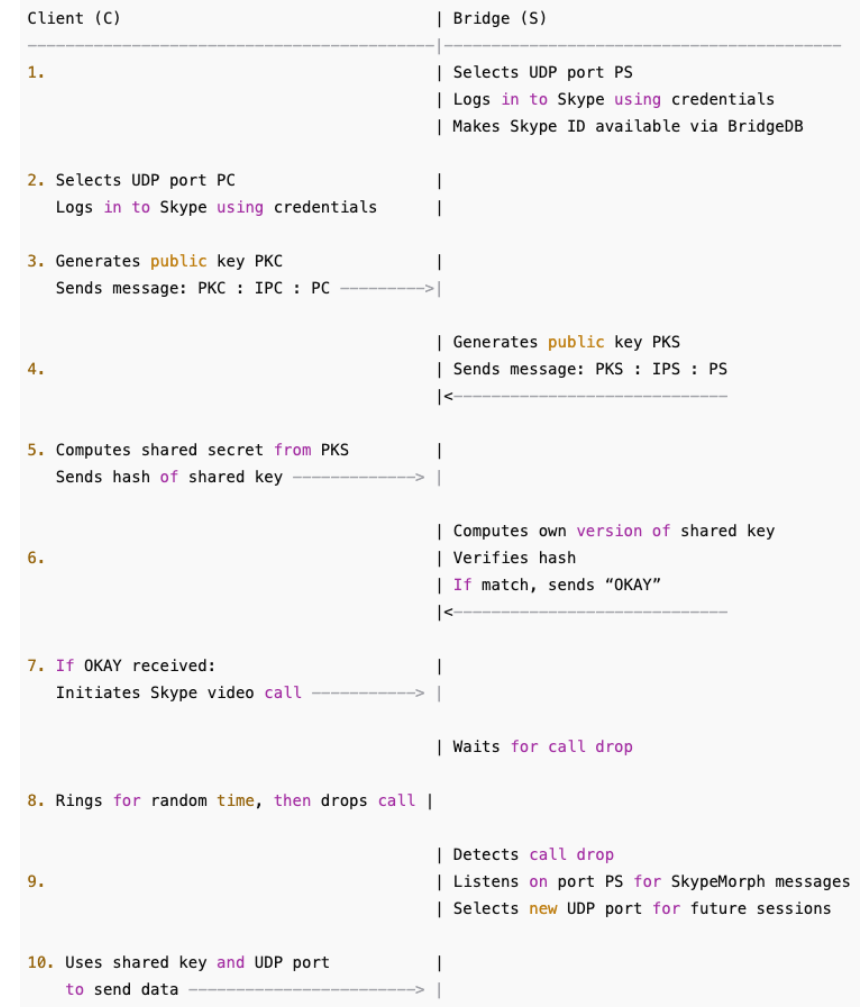
Conversation setup:

Protocol between client and bridge using the Skype API to establish conversation

- Use UDP protocol as vanilla Skype
- Initiate call that is then dropped
- Use agreed UDP port to exchange packets

→ Now can send packets to each other.

Is there anything else client and bridge need to take into account?



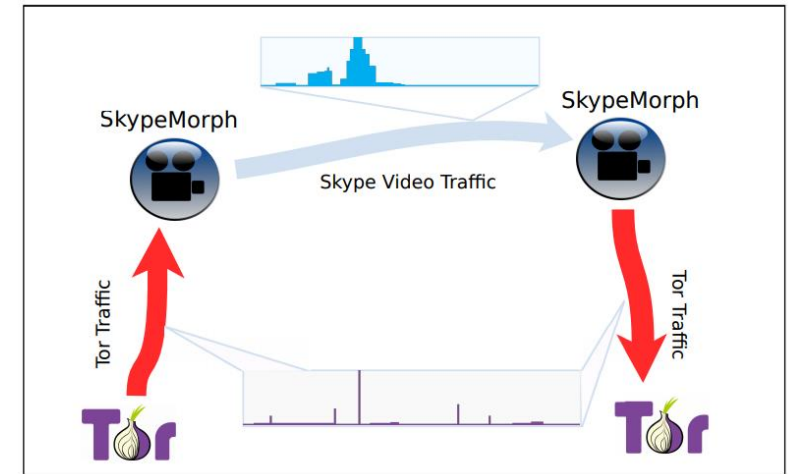
SkypeMorph

Traffic shaping

Exchange: Send Tor TLS data over encrypted channel, masquerading it as Skype video

- Mimic skype traffic
 - Packet size
 - Inter-arrival times

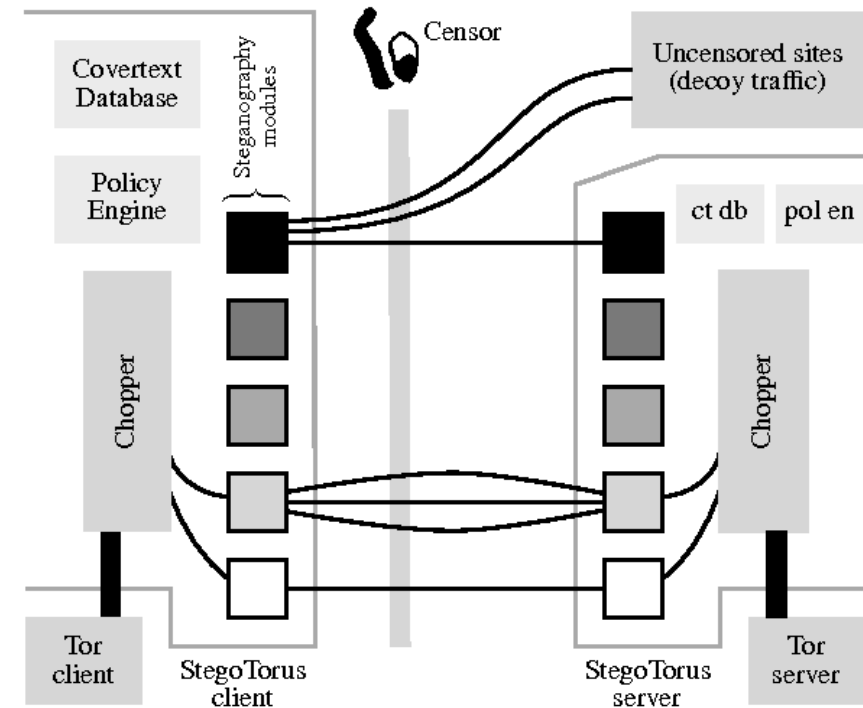
→ **Packets are sent (statistically) following Skype patterns**



Mimicry: Look like whitelisted StegoTorus

Goal: Make it difficult for the censor to distinguish between the obfuscated bridge connections and whitelisted traffic using statistical comparisons of flow features

How: Chops Tor traffic and sends it through different connections (HTTP, Skype, VoIP...)

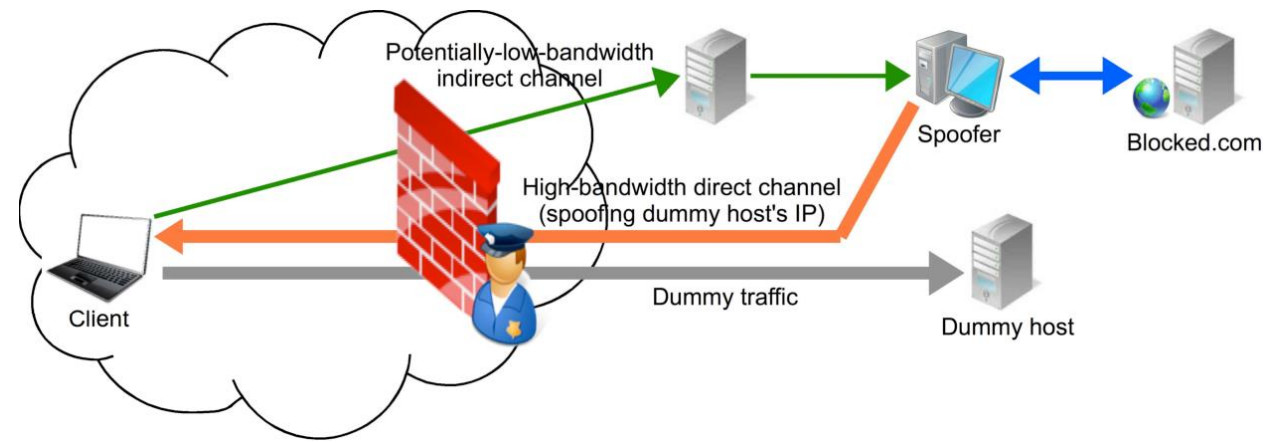


Mimicry: Look like whitelisted CensorSpoofer

Goal: Obfuscate traffic patterns through mimicry

How: Standalone system

- (1) IP Spoofing to obfuscate server's identity
- (2) Mimics VoIP traffic to obfuscate traffic patterns



Mimicry: Look like whitelisted

Goal: Make it difficult for the censor to distinguish between black- and whitelisted traffic using statistical comparisons of flow features

How: Imitate common protocols like HTTP and Skype

Systems: SkypeMorph, StegoTorus, CensorSpoof...

→ Parrot circumvention systems



Mimicry: Look like whitelisted

The parrot is dead

Goal: Make it difficult for the censor to distinguish between black- and whitelisted traffic using statistical comparisons of flow features

How: Imitate common protocols like HTTP and Skype

Systems: SkypeMorph, StegoTorus, CensorSpoofer...



→ “Unobservability by imitation” is fundamentally flawed.

Censorship

Step 1: Fingerprinting

Destination:  **Tor (other anon comms)**
IP addresses, hosts, ports,...

Content:  **Encryption**
protocol strings, keywords, domains, http hosts, encrypted flows...

Flow properties:  **Obfuscation through mimicry**
length, inter-arrival times, bursts, ...

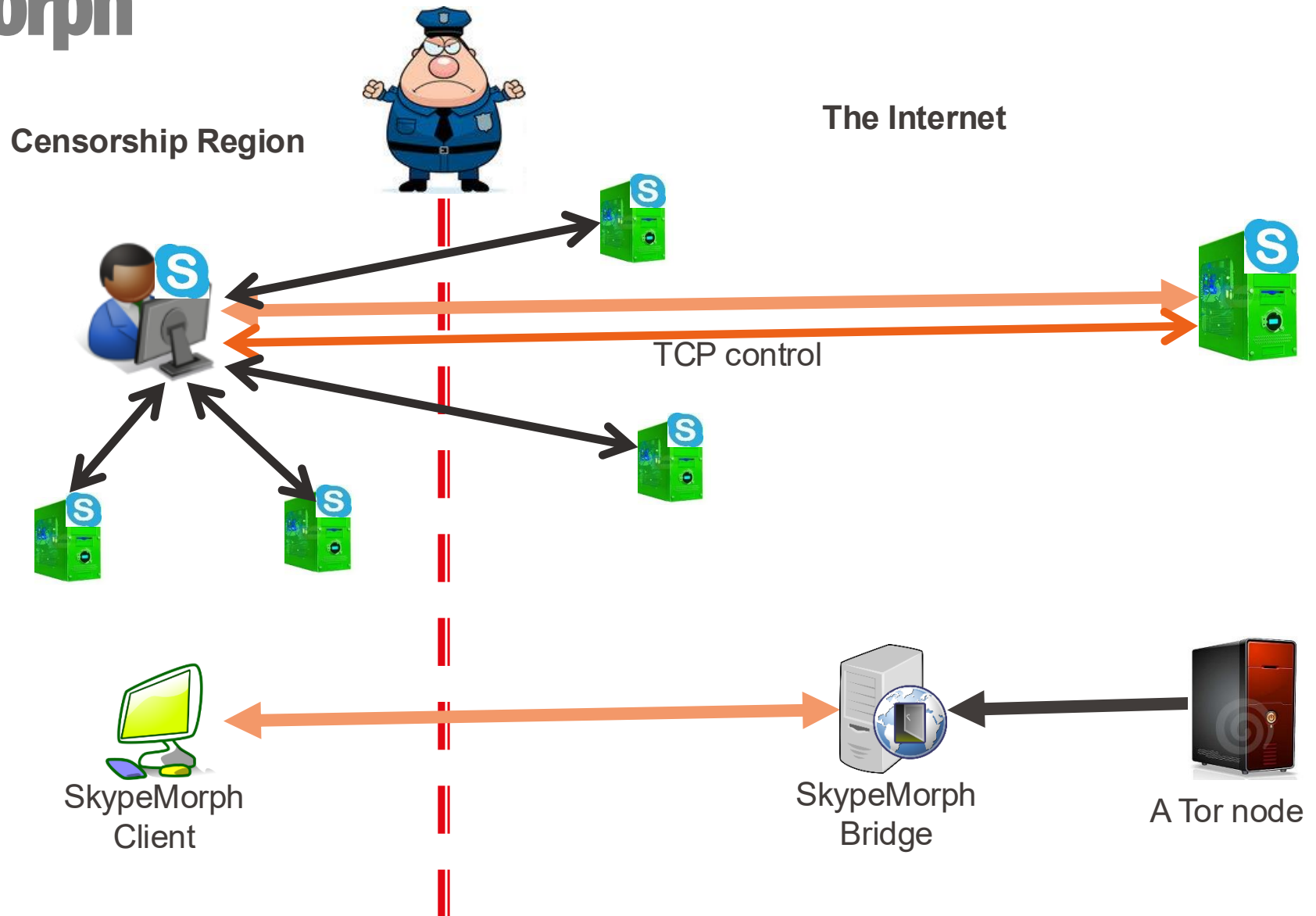
Protocol semantics:
protocol behavior (mostly active attacks)

To win, the censor needs only to find a few discrepancies



The parrot is dead

SkypeMorph

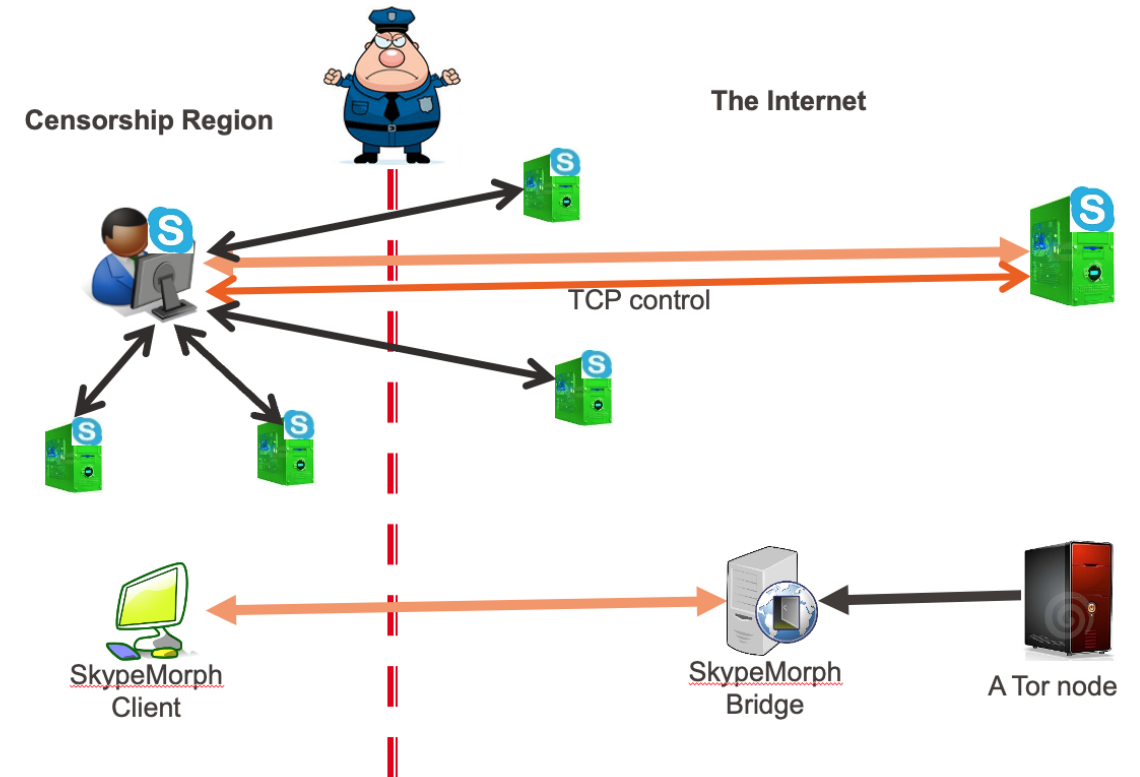


Parrots mimic Skype's traffic statistics but...

...fail to mimic much more visible aspects:

- no HTTP update traffic
- no login traffic
- no mimicry of Skype's TCP channel

→ Parrot systems can be distinguished from Skype even by extremely basic tests



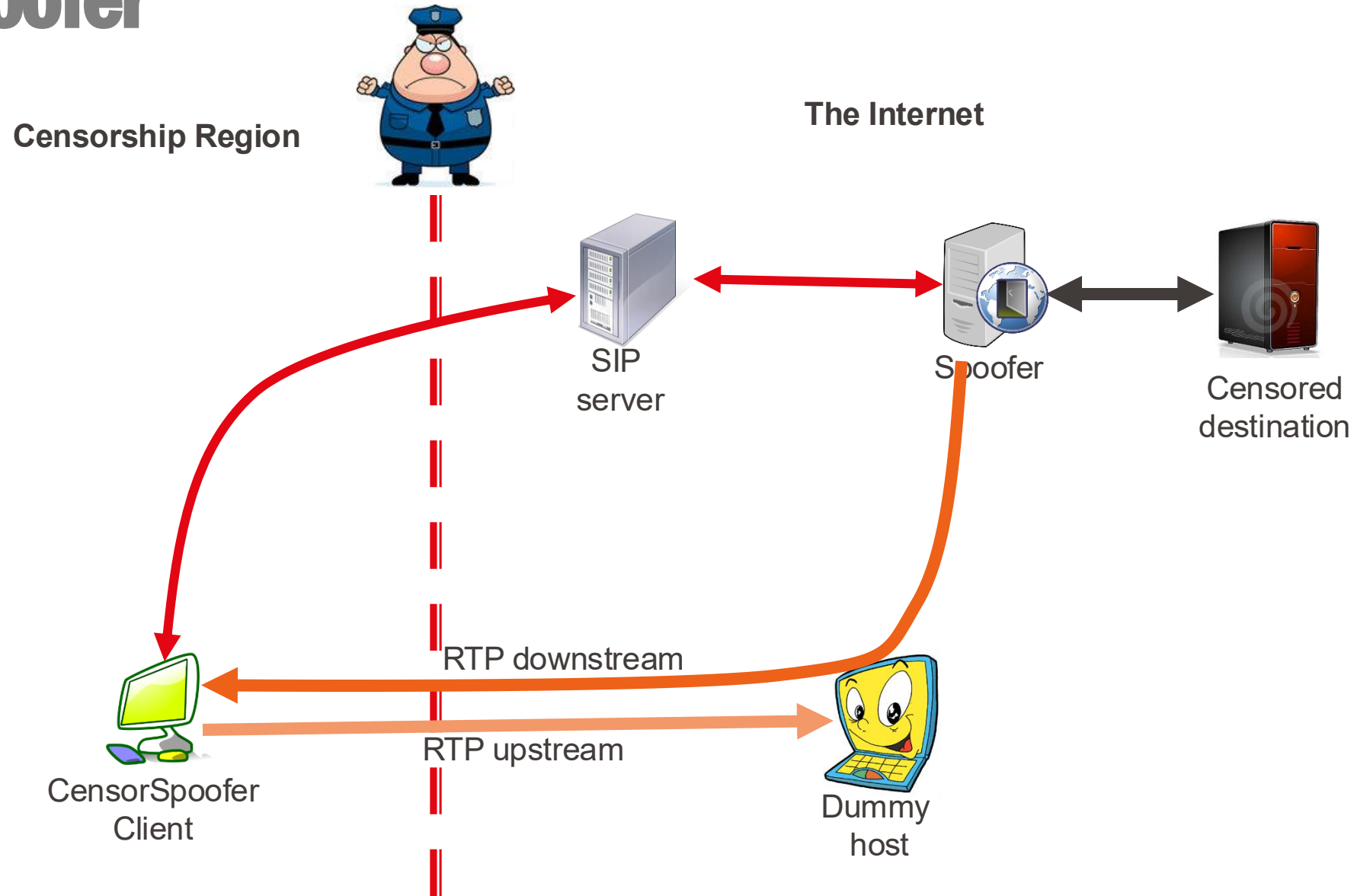
The parrot is dead

Other tests

Test	Skype	SkypeMorph+
Flush Supernode cache	Serves as a SN	Rejects all Skype messages
Drop UDP packets	Burst of packets in TCP control	No reaction
Close TCP channel	Ends the UDP stream	No reaction
Delay TCP packets	Reacts depending on the type of message	No reaction
Close TCP connection to a SN	Initiates UDP probes	No reaction
Block the default TCP port	Connects to TCP ports 80 and 443	No reaction

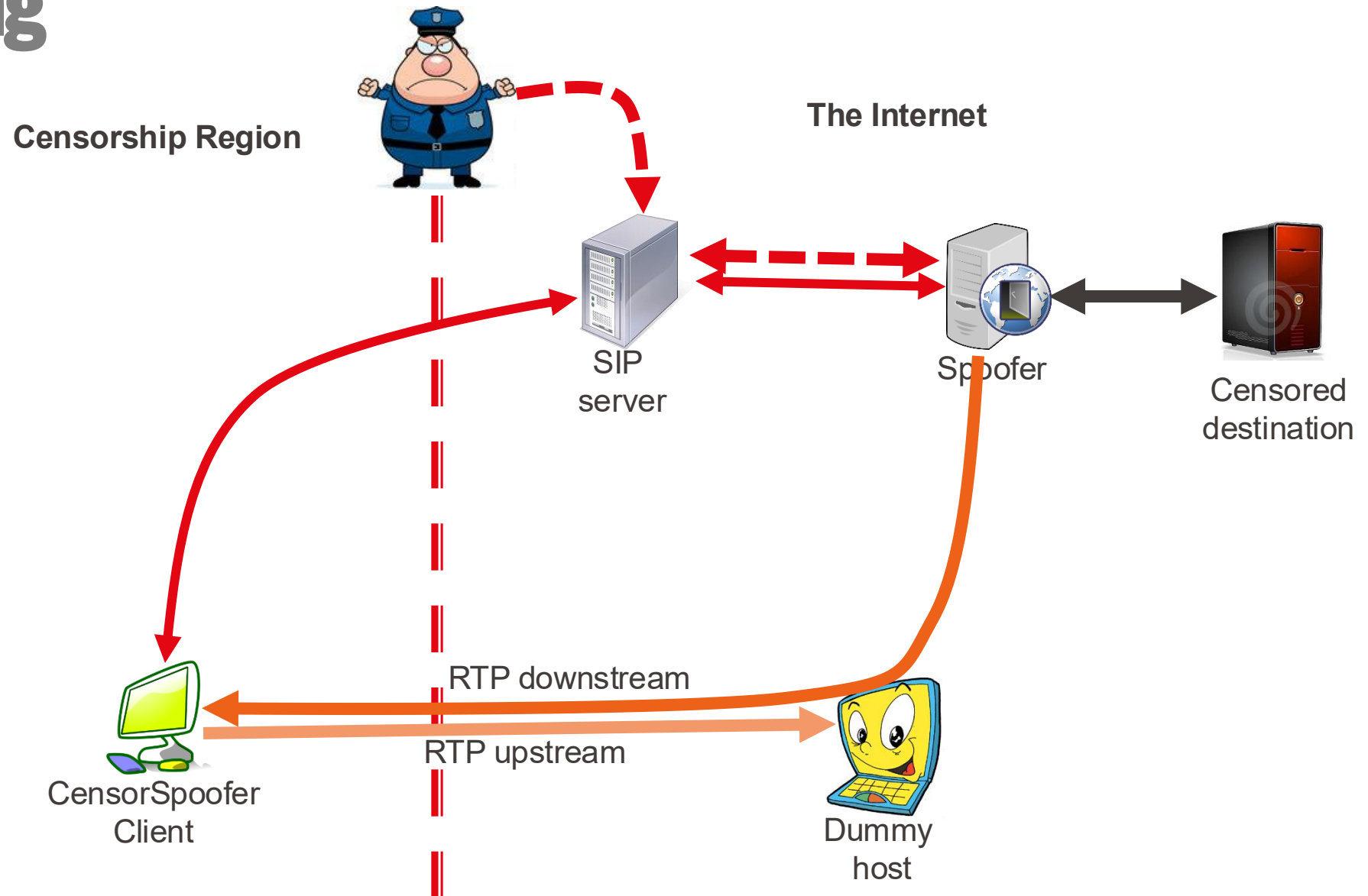
The parrot is dead

CensorSpoofer



The parrot is dead

SIP probing



The parrot is dead

Imitation Requirements

Parrot needs to mimic...

Protocol in its entirety	Reaction to errors and network conditions	Typical traffic
Correct protocol	Errors	Content
Side protocols	Network	Patterns
Intra dependencies		Users
Inter dependencies		Geolocalisation

Solution: do not imitate, be!!



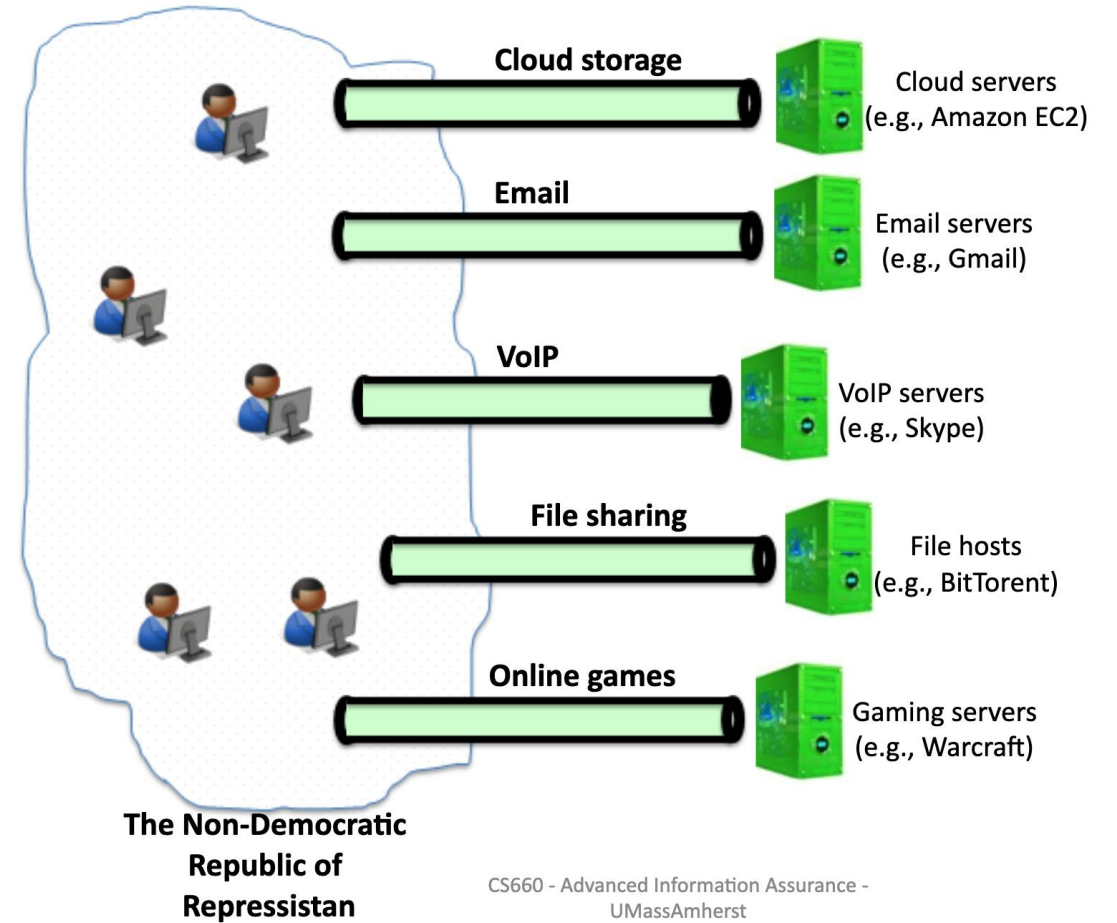
From parrots...



...to parasites

Idea: We already have a lot of encrypted channels...

→ Hide censored traffic within!



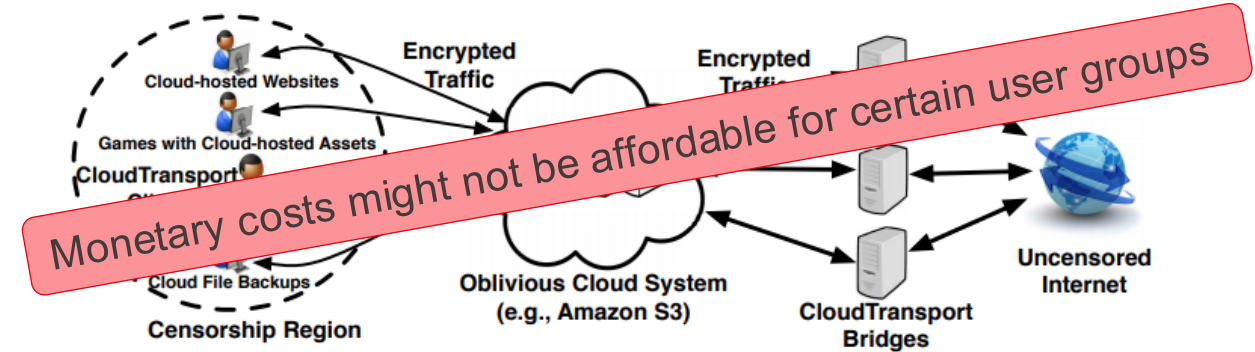
Hide-within circumvention

CloudTransport

Goal: raise economic and social costs of censorship by forcing the censors to use statistical traffic analysis and other computationally intensive techniques

How: Hide censored traffic within existing encrypted channels

- 1) Select a Cloud provider: one that does provide other non-censored services
- 2) Create a rendez-vous account with the Cloud ["can't" be censored!]
- 3) Select a CloudTransport bridge and send to it the rendez-vous credentials (Dead drop or Out-of-band)
- 4) To send data, the client puts it on the Cloud and the Bridge transmits it to the destination



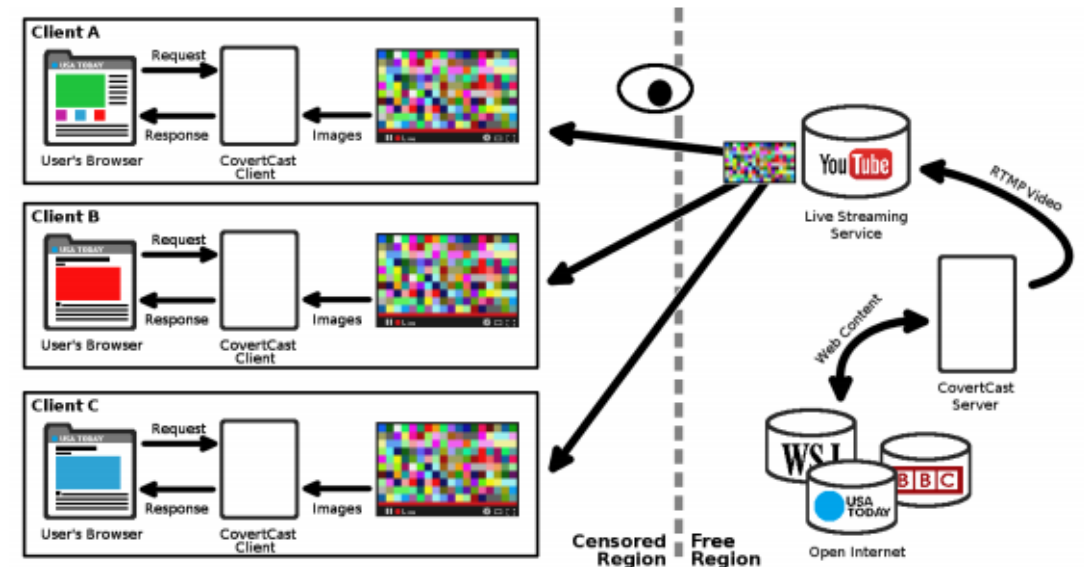
Hide-within circumvention

CovertCast

Goal: raise economic and social costs of censorship by forcing the censors to use statistical traffic analysis and other computationally intensive techniques

How: Hide censored traffic within existing encrypted channels

- 1) CovertCast server initiates live stream
- 2) Server crawls censored site, encodes content into images and broadcasts images via live stream
- 3) CovertCast client demodulates images back into Web content



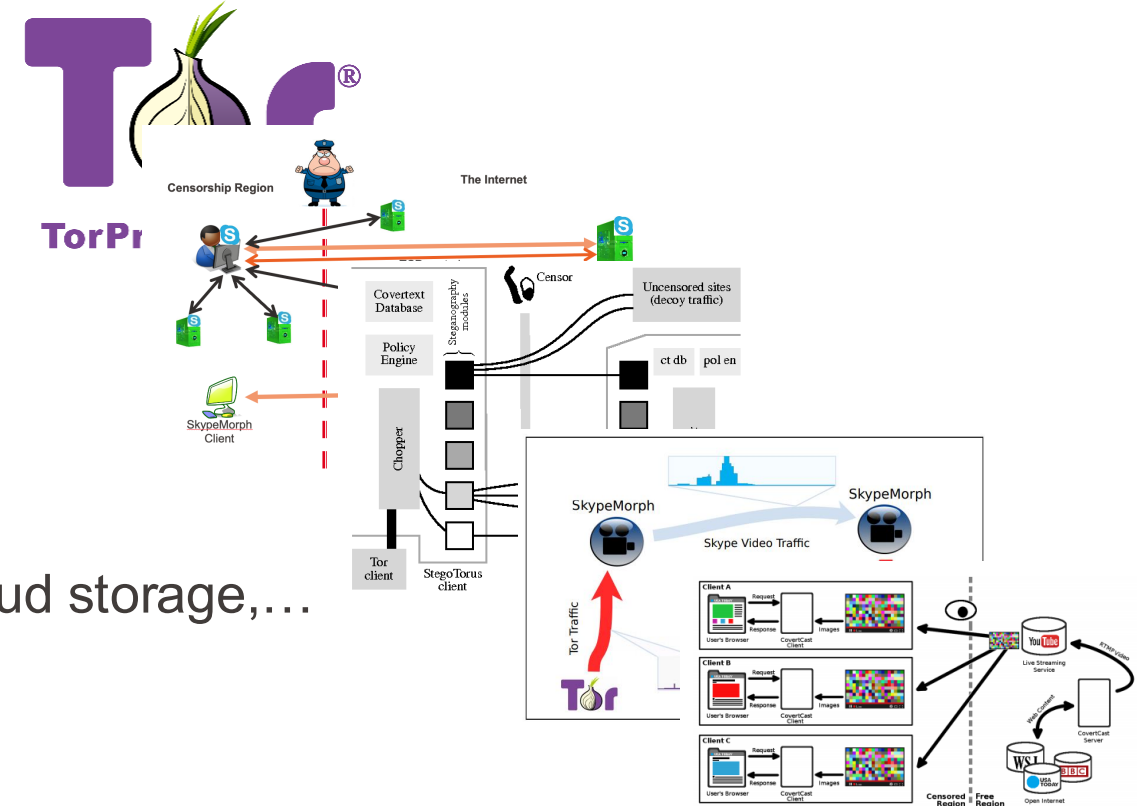
All systems work **at the application layer**:

Overlay networks: Onion routing & obfuscation

Reuse other infrastructures:

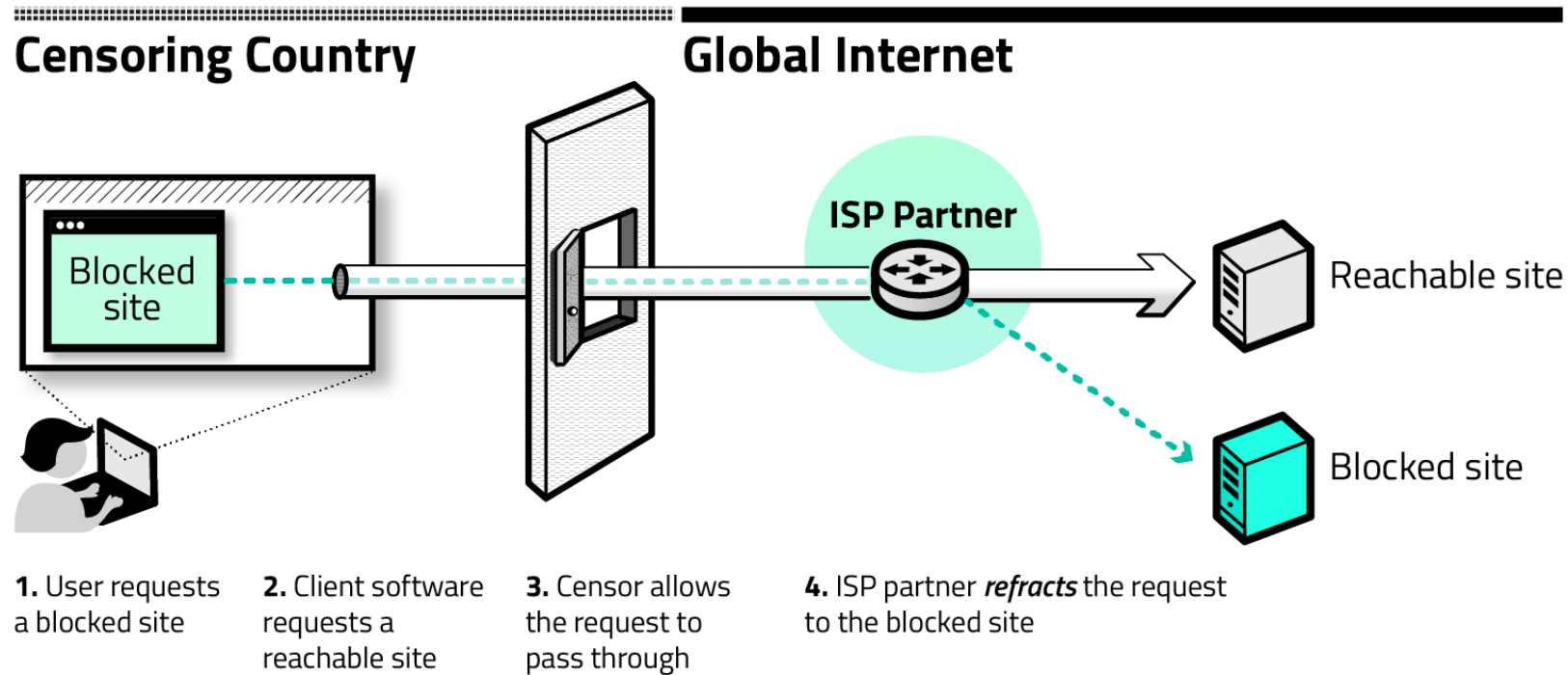
Parrots: Imitation of Skype, P2P,...

Hide-within: Hide within live streams, cloud storage,...



→ End the cat-and-mouse game of application-layer censorship systems

Motivation: End the cat-and-mouse game of application-layer censorship systems

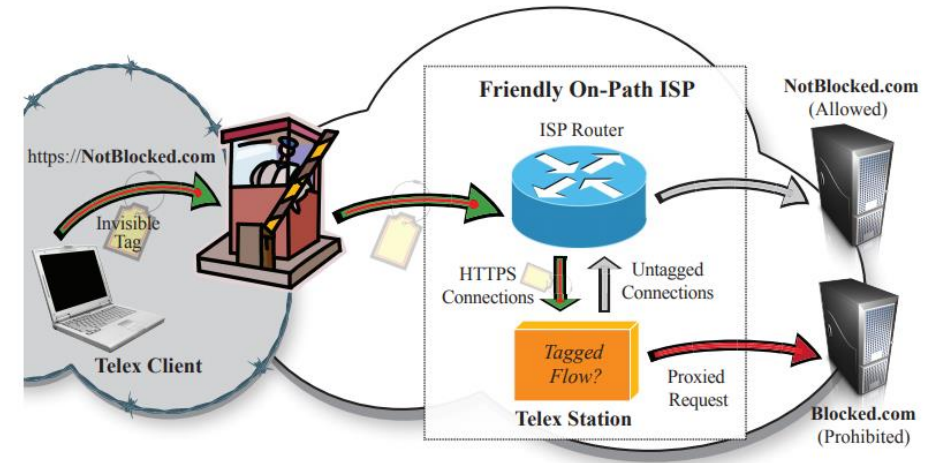


<https://refraction.network>

Decoy routing

Example: Telex

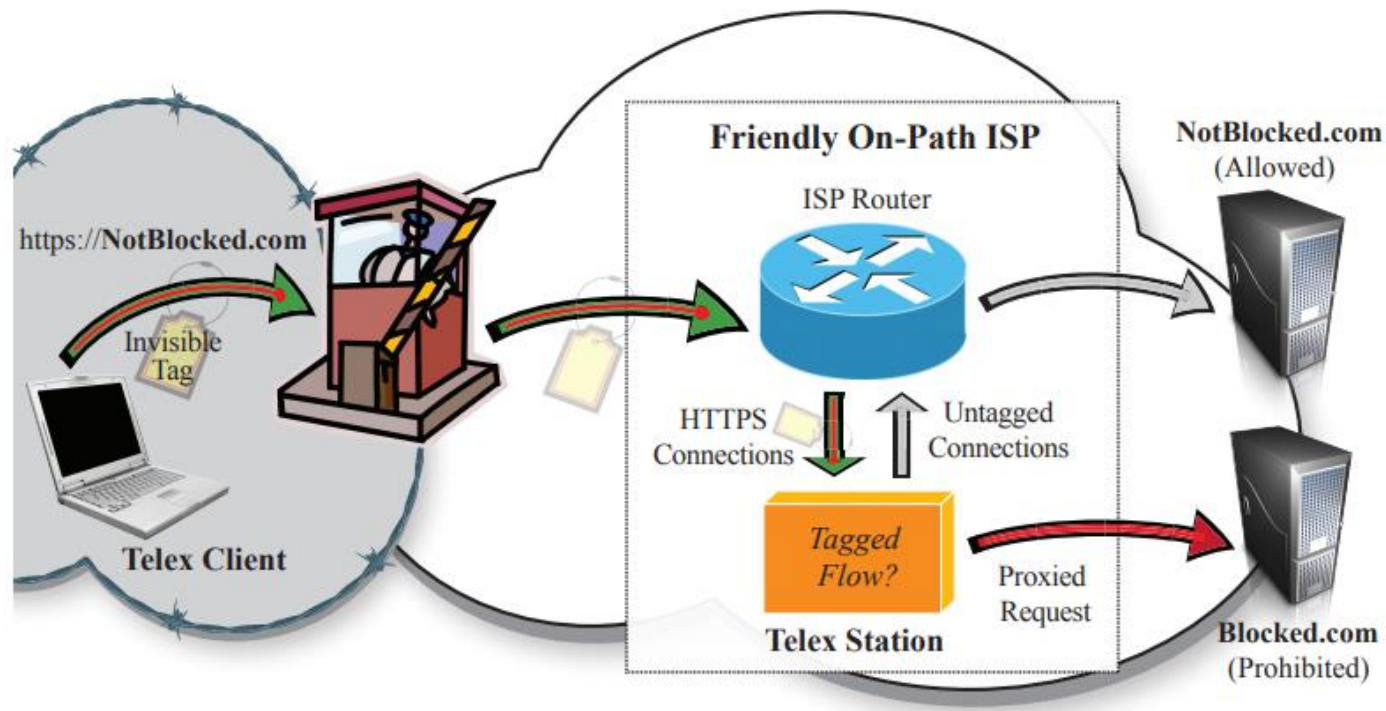
- Operates in the **network infrastructure** — at any ISP between the censor's network and non-blocked
 - **State-level response** to state-level censorship.
- Repurposes **deep-packet inspection** to circumvent censorship.
- **No secrets** to communicate to users in advance
- Focuses on **avoiding detection**



“A friendly man-in-the-middle”

Decoy routing

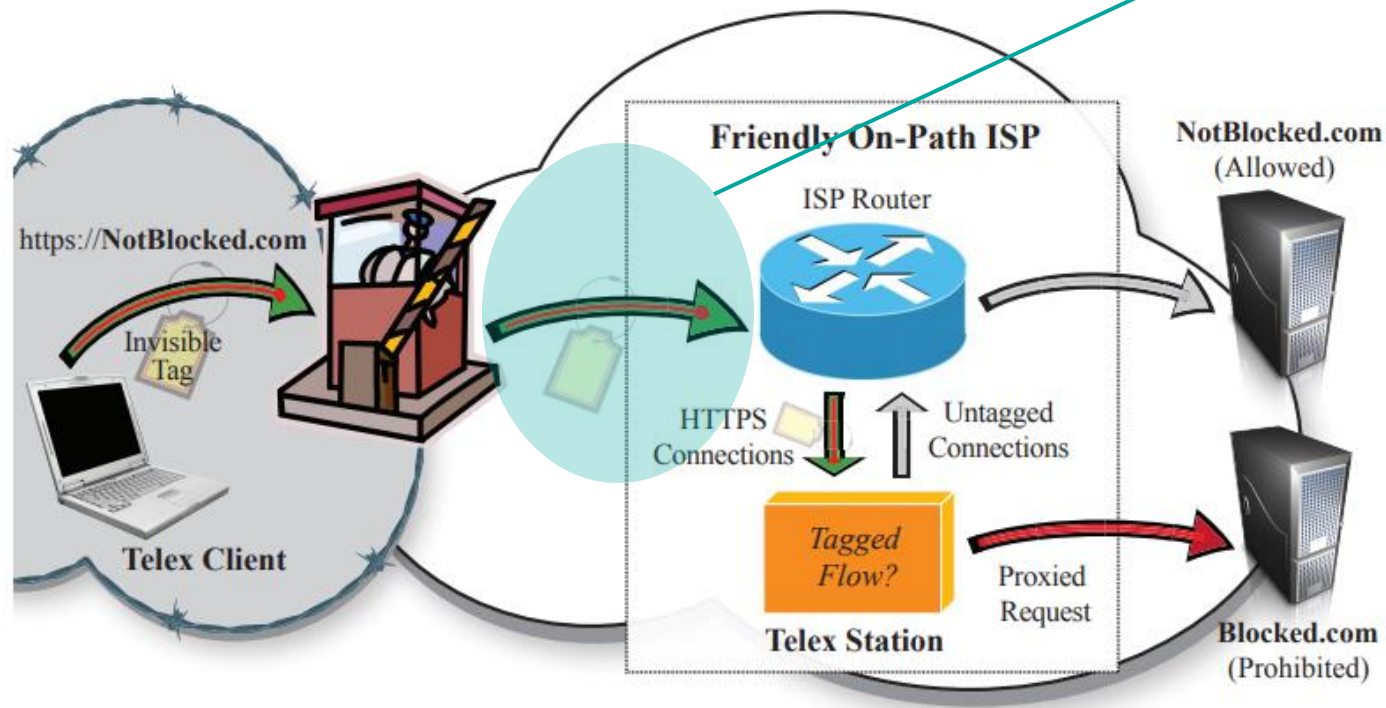
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Decoy routing

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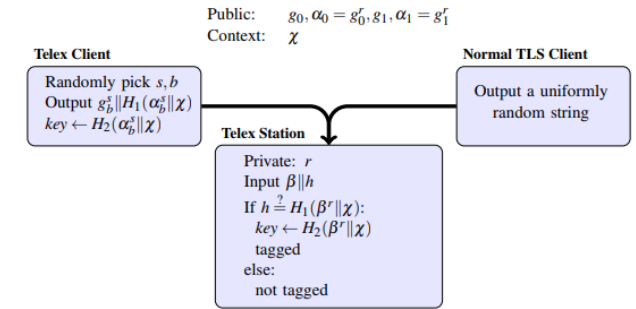
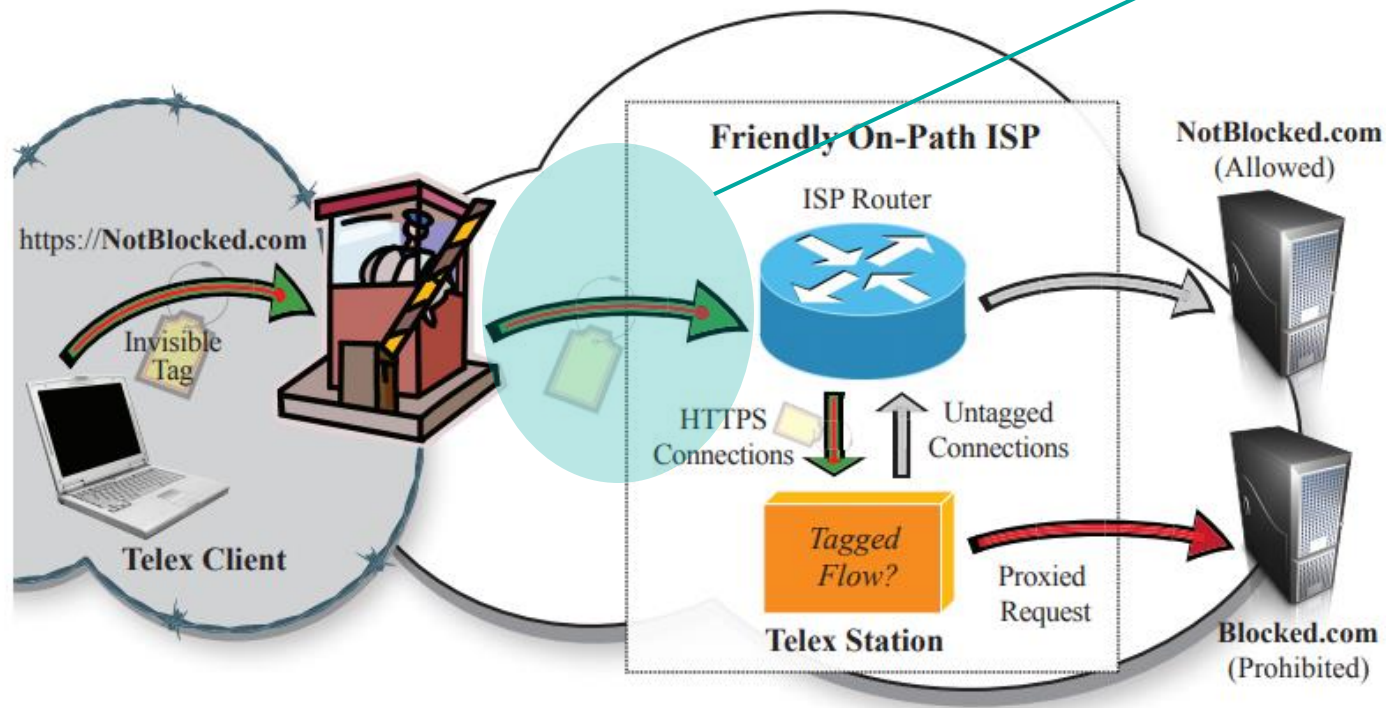
Tag: looks like a random nonce in the TLS handshake



Decoy routing

Example: Telex

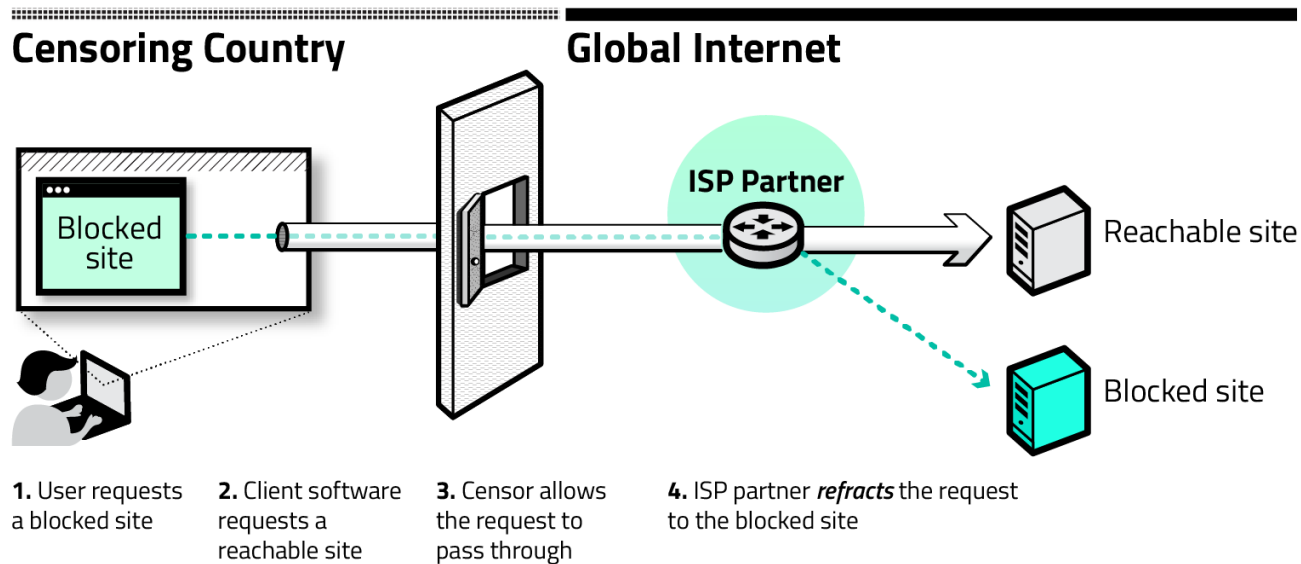
Tag: looks like a random nonce in the TLS handshake



Decoy routing

Routing attacks

Motivation: End the cat-and-mouse game of application-layer censorship systems



<https://refraction.network>

Have we really reached the end?

Meet routing capable adversaries



A censoring authority who is capable of controlling how packets originating from its network are routed

Decoy routing

Routing attacks

Through routing attacks a routing capable adversary can:

- Enumerate the participating decoy routers
- Successfully avoid sending traffic along routes containing these routers with little or no adverse effects
- Identify users of these schemes through active and passive attacks
- (In some cases) probabilistically identify connections to targeted destinations.

One of the goals of Telex:
Avoid detection



A censoring authority who is capable of controlling how packets originating from its network are routed

Decoy routing

Routing attacks

How: Routing adversary must be able to

1) Locate decoy routers

Telex: Public list of decoy router locations

Ciripede: Scan Autonomous Systems (ASs)

→ Make a list of honest vs. tainted ASs

2) Select from a diverse set of paths in reaction to this knowledge

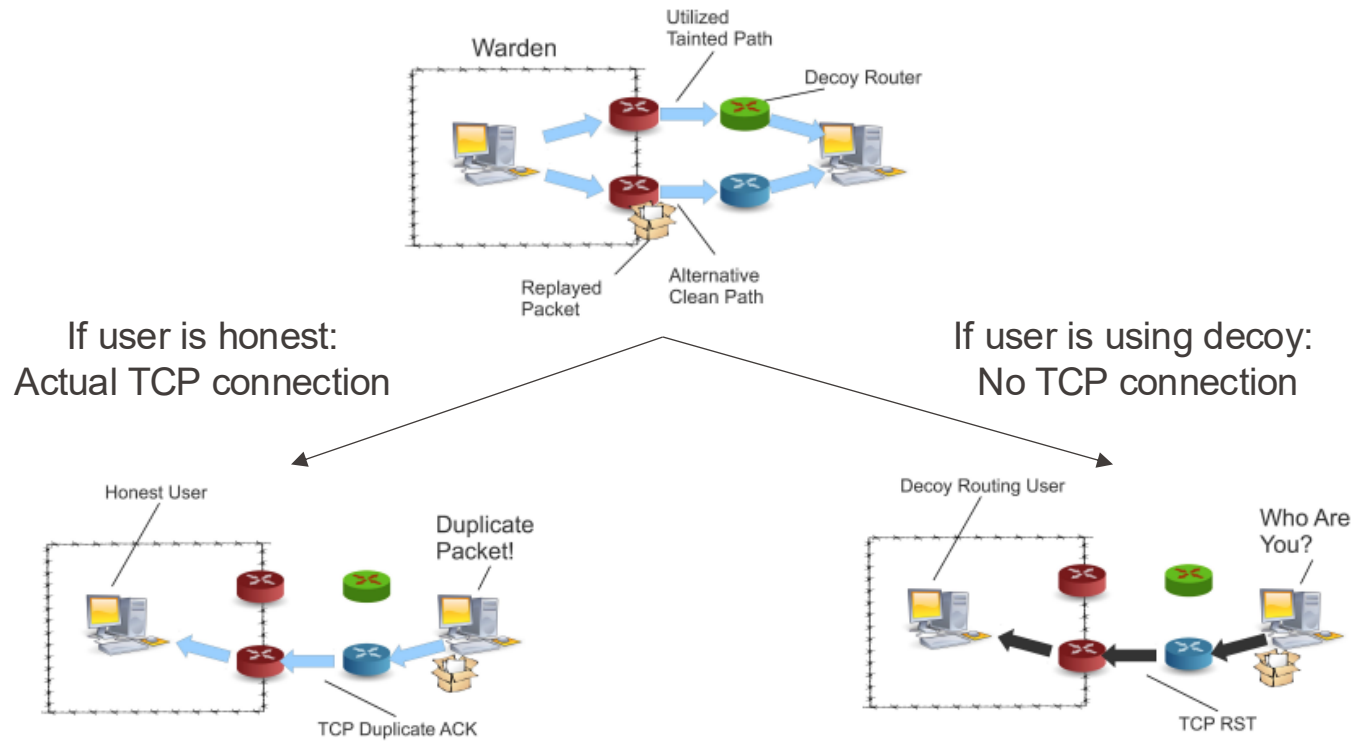


A censoring authority who is capable of controlling how packets originating from its network are routed

Routing attacks

Detection attacks

TCP Replay attack

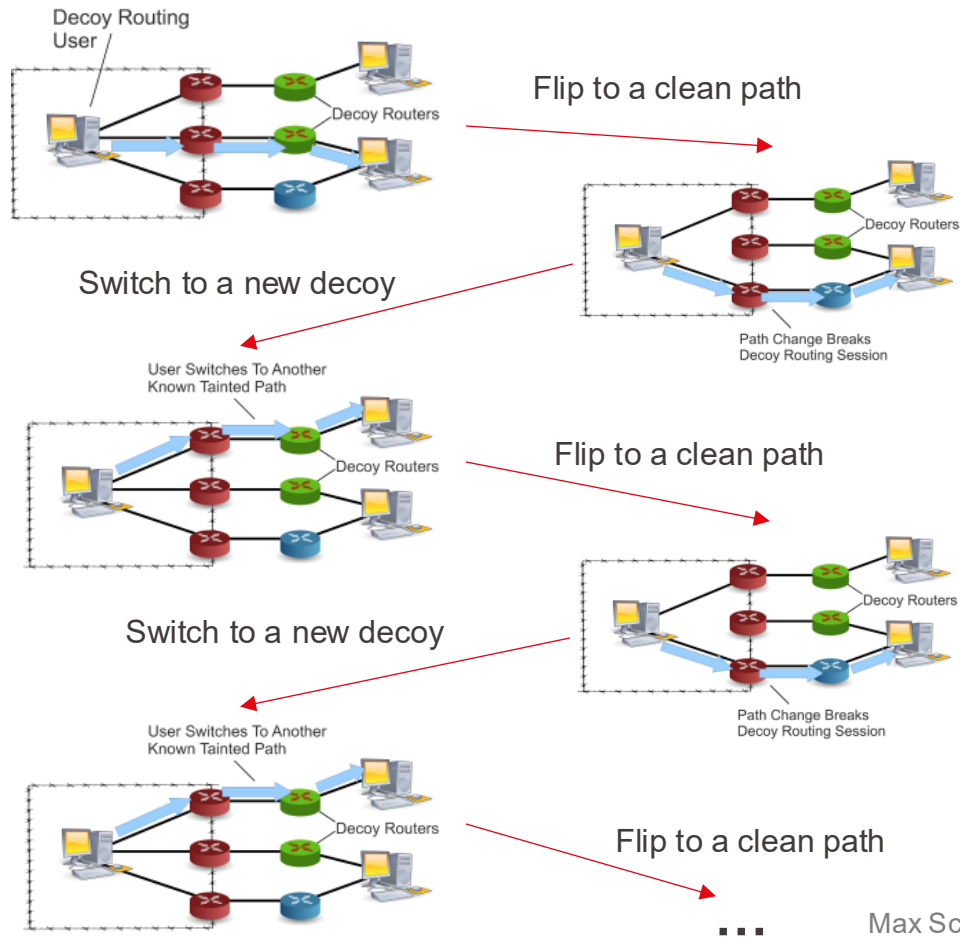


Goal: Identify users of decoy routing systems

Routing attacks

Detection attacks

The “Crazy Ivan” attack



Goal: Identify users of decoy routing systems

Routing attacks

Timing attacks

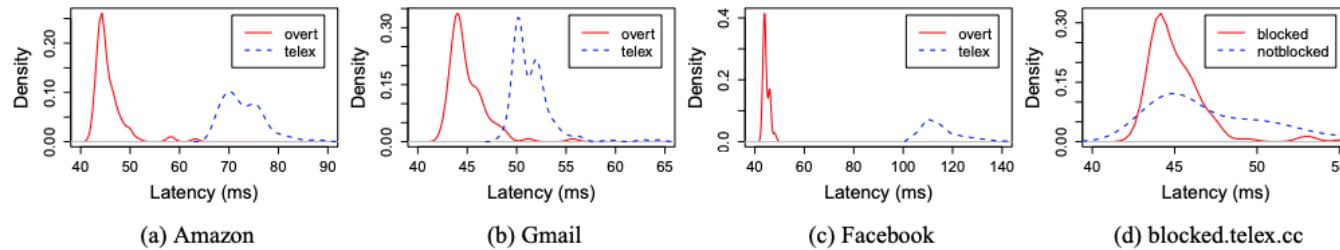


Figure 5: Comparing distribution of latencies from notblocked.telex.cc to (a) Amazon (b) Gmail (c) Facebook and (d) blocked.telex.cc

How: Fingerprint network latency



Goal: Identify users of decoy routing systems

Decoy routing

Routing attacks

Through **routing attacks** a routing capable adversary can:

- Enumerate the participating decoy routers
- Successfully avoid sending traffic along routes containing these routers with little or no adverse effects
- Identify users of these schemes through active and passive attacks
- (In some cases) probabilistically identify connections to targeted destinations.

CRAZY EXPENSIVE!!



A censoring authority who is capable of controlling how packets originating from its network are routed

Take aways

- Censorship resistance is key to freedom speech & information
- There is a strong connection between censorship resistance technology and anonymous communications
- To resist internet censorship requires:
 - Bootstrapping: find “helper” nodes
 - Lists, private retrieval, embedded in infrastructure
 - Hidden communication: avoid censor “during conversation”
 - Hide: network information, content, patterns
 - Comply with semantics ← do not imitate, be